Augmenting Batch Exchanges with Constant Function Market Makers EC 2024

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Outline

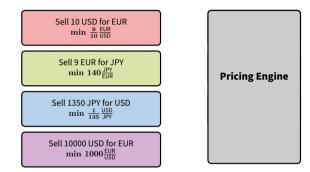
- Two ideas in exchange design with newfound popularity
- How should we combine them?
- Goal: Map out design space (no dominant design)

Exchange Model

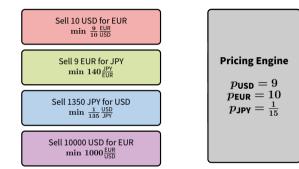
• Users trade *N* divisible, fungible assets through *limit orders*

- "Sell 1 unit of $\mathcal X$ for at least 2 units of $\mathcal Y$ "

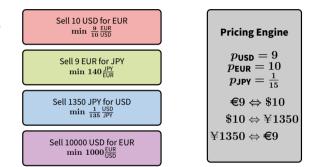
- Execute batches of trades, all at once
- Input: Set of limit orders
- 1. Compute Prices
- 2. Trade in batch at price quotients
 - Meaningless units
 - No pairwise matching
 - "Clearing" if no debt



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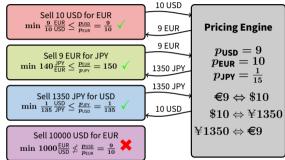


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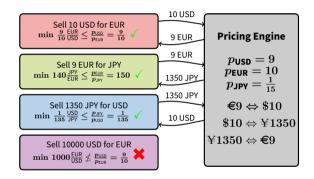
Theorem (Arrow and Debreu, 1954)

 \exists unique^{*} equilibrium prices $\{p_A\}$ and allocations that clear the market.



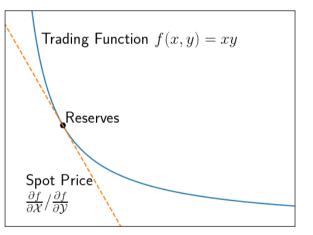
Key Properties of Batch Exchange Model

- **1** Uniform prices (unique!) bring economic benefits
 - Pareto-Optimal (for limit orders)
 - E.g. Budish et al. "The high-frequency trading arms race" (2015)
- 2 Requires computing Arrow-Debreu exchange market equilibria



Two Exchange Design Innovations Constant Function Market Makers

- CFMM maintains reserves and a trading function f(·)
- Accepts trade from (x, y) to (x', y') if and only if $f(x, y) \le f(x', y')$
- Why?
 - Market-makers add liquidity
 - Automated
 - Computational simplicity



Our Work In Context

- Several projects combine batch exchanges with CFMMs, using different mechanisms
 - Penumbra, CoWSwap, [Walther, 2021], [Canidio and Fritsch, 2023]
- What are the tradeoffs for different mechanisms for integrating CFMMs into batch exchanges?

Augmenting Batch Exchanges with CFMMs

How can batch exchanges draw on passive liquidity?

- Model:
 - N assets $\mathcal{X} \in \mathfrak{A}$
 - 1 batch exchange
 - Many CFMMs, with different curves, reserves
 - Also outside world—other exchanges, other users, ...

Augmenting Batch Exchanges with CFMMs

How can batch exchanges draw on passive liquidity?

- Axiom 1: Asset conservation
- Axiom 2: Uniform Prices $\{p_{\mathcal{X}}\}_{\mathcal{X}\in\mathfrak{A}}$
 - No trade from \mathcal{X} to \mathcal{Y} gets a better rate than $\frac{p_{\mathcal{X}}}{p_{\mathcal{Y}}}$.

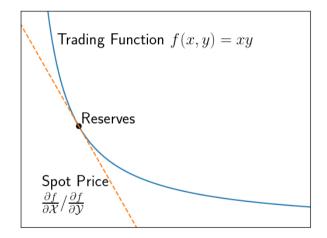
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- Axiom 3: Limit orders make best responses
 - A limit order trades \mathcal{X} to \mathcal{Y} at no worse than the market rate $\frac{p_{\mathcal{X}}}{p_{\mathcal{Y}}}$, only if market rate exceeds limit price
- These are standard market design assumptions, lead to classic theory results on Arrow-Debreu market equilibria.

How can batch exchanges draw on passive liquidity?

 Axiom 4: CFMM trading function must not decrease

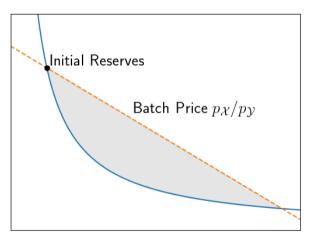


How can batch exchanges draw on passive liquidity?

 Axiom 4: CFMM trading function must not decrease

Consequence

Market equilibrium is no longer unique



CFMMs in Batch Exchanges

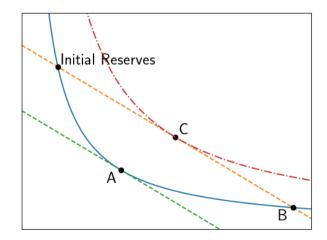
How can batch exchanges draw on passive liquidity?

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Consequence

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- How should a batch choose a CFMM's trade?
- Also complicates equilibria computation

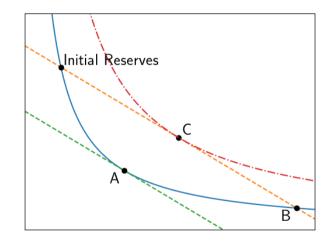


CFMMs in Batch Exchanges

Asset Conservation and Uniform Prices imply:

Consequence

CFMMs must trade *at* market prices, not below



Some Desirable Properties

Pareto Optimality

- From perspective of limit orders
- Recall: Without CFMMs, every equilibrium is Pareto Optimal

Price Coherence

- After a batch, CFMM spot exchange rates are quotients of some set of prices
- Otherwise, cyclic arbitrage opportunity (free money)

Preservation of Price Coherence

- Price coherence, but only if prices are also coherent before batch

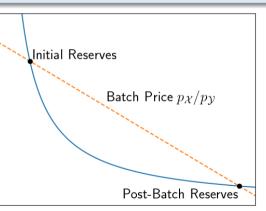
Consequences

Consequence

No mechanism can in all circumstances guarantee Pareto Optimality and (Preservation of) Price Coherence

• Proof Intuition:

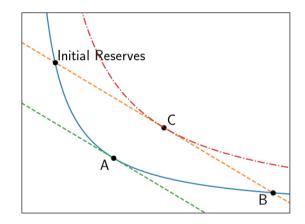
- PO can require trading all the way across
- Multiple CFMMs with different curves will end at different spot exchange rates



Some More Desiderata and Consequences

Joint Price Discovery (JPD)

- After a batch, CFMM spot prices equal batch prices
- Prevents a common atomic, risk-free "cyclic" arbitrage
- JPD requires maximizing f(·) (trading to C)
- An example of how context matters:
 - Trading to *C* incentivizes splitting trade over many batches, but trading to *B* does not.
 - How are batches initiated?
 - How many users?



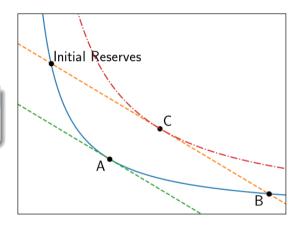
Some More Desiderata and Consequences

• Locally Computable Rule (LCR)

- CFMM trade depends only on CFMM state and market price

Consequence

Trading to C is a LCR that satisfies Price Coherence

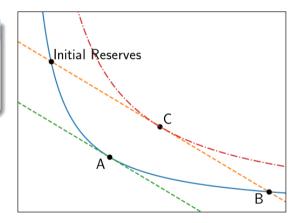


Some More Desiderata and Consequences

Consequence

Trading to *B* is a LCR that guarantees Preservation of Price Coherence, if and only if all CFMMs use a *constant product* curve.

• Unique exception to incompatibility between PO and PPC



Computing Equilibria

- Mixed-Integer Programs [Walther21] or general (not always convex) solvers
- LCR \Rightarrow algorithms based on auctions, iterations (Tâtonnement) are directly applicable
 - LCR must satisfy Weak Gross Substitutability
 - Price goes up \implies demand does not increase
- What about other approaches? Convex programs?

A Convex Program for 2-Asset WGS Utility Functions

Observation

A CFMM trading between 2 assets, with a LCR satisfying WGS, acts like an (uncountably) infinite set of infinitesimal limit orders.

Let's adapt a convex program for linear Arrow-Debreu exchange markets [DGV16] to support CFMMs trading between 2 assets

A Convex Program for Linear Utility Functions [DGV16]

$$\begin{array}{ll} \text{Minimize } \sum_{i} p_i \left(e_i \, \ln(\frac{p_i}{\beta_i}) \right) - \sum_{i} y_{i,j} \ln u_{i,j} \\ \text{Subject to } \sum_{i} y_{i,j} = \sum_{i} y_{j,i} & \forall j \in [N] \\ p_j \geq 1 & \forall j \in [N] \\ y_i \geq 0 & \forall i \in [M] \\ u_{i,j} \beta_i \leq p_j & \forall i,j \end{array}$$

A Convex Program for 2-asset WGS CFMM Trading Functions

$$\begin{array}{ll} \text{Minimize } \sum_{i} p_{A_{i}} \int_{0}^{\infty} \left(d_{i}(z) \ln(\frac{p_{A_{i}}}{\beta_{i,z}(p)}) \right) dz - \sum_{i} p_{A_{i}} g_{i}(y_{i}/p_{A_{i}}) \\ \text{Subject to } \sum_{i:A_{i}=j} y_{i} = \sum_{i:B_{i}=j} y_{i} \\ p_{j} \geq 1 \\ y_{i} \geq 0 \end{array} \qquad \forall j \in [N] \\ \forall i \in [M]. \end{array}$$

Equivalently, this program solves exchange markets where each agent is interested in only two assets, using *any* WGS utility function on those two assets.



• Axiomatic framework for integrating CFMMs into batch exchanges

- Extra degree of freedom requires deliberate choice
- Natural desiderata are incompatible
 - Pareto-Optimality at odds with Price Coherence
- Convex program for exchange markets with 2-asset WGS CFMMs